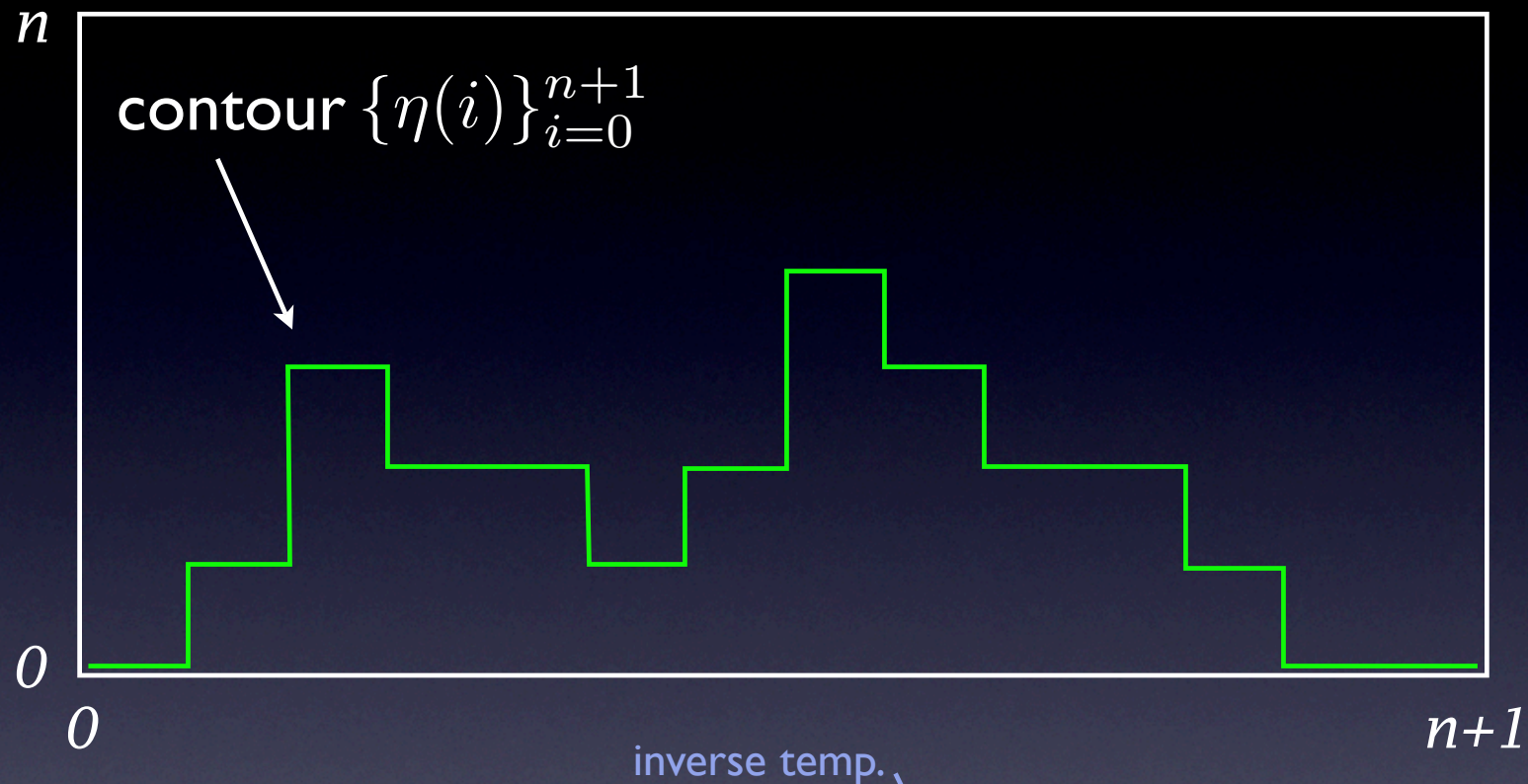


Mixing time for the Solid-on-Solid model

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joint work with
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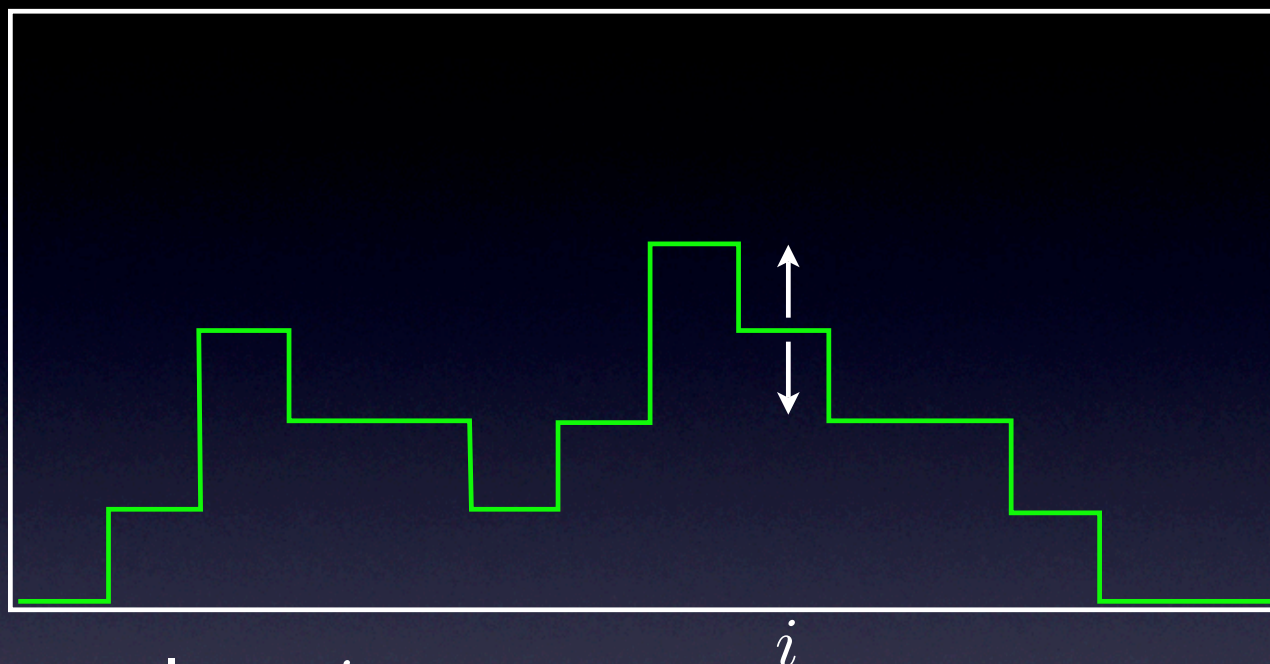
Solid-on-Solid (SOS) model



Gibbs distribn: $\pi(\eta) = Z_{\beta}^{-1} \exp\left\{-\beta \sum_{i=1}^{n+1} |\eta(i) - \eta(i-1)|\right\}$

Boundary conditions: $\eta(0) = \eta(n+1) = 0$

Single-site (Glauber) dynamics



- pick a column i u.a.r.
- set $\eta(i) \in \{\eta(i), \eta(i) \pm 1\}$ with appropriate probs.
- $\eta_t \rightarrow \pi$ as $t \rightarrow \infty$

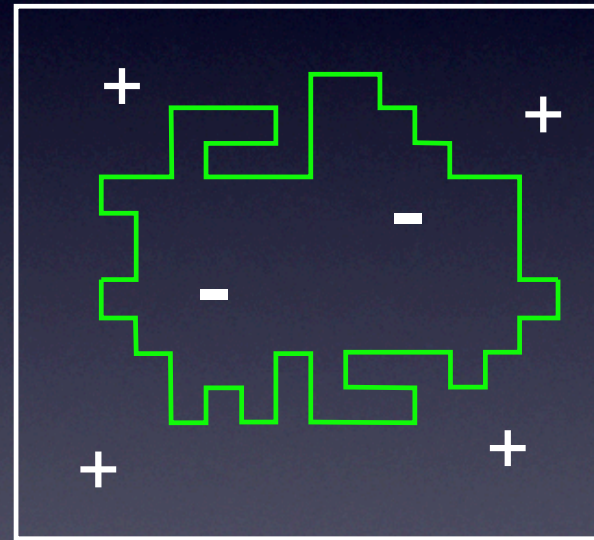
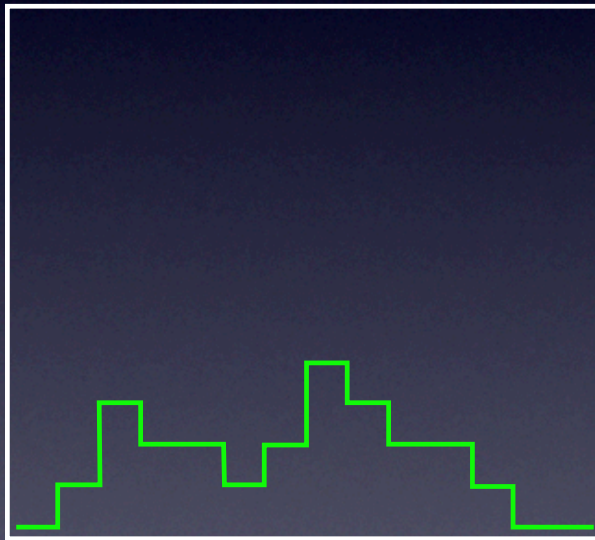
Goal: analyze the mixing time:

$$\tau = \min\{t : \|\eta_t - \pi\| \leq 1/4 \ \forall \eta_0\}$$

Why?

1. Model for random surfaces etc. [Privman/Švrakić...]

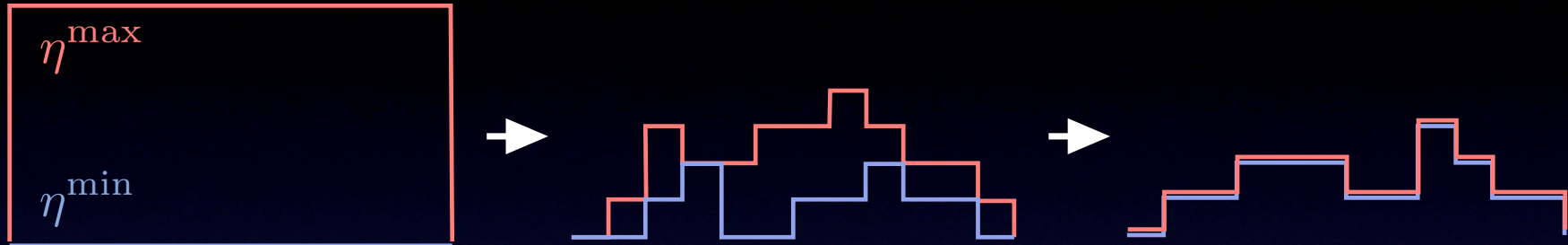
2. Connection with low temperature Ising model



At low temps, few “overhangs” \Rightarrow good approximation
(Zero temp. solved by [Chayes/Schonmann/Swindle])

3. Challenge to existing techniques

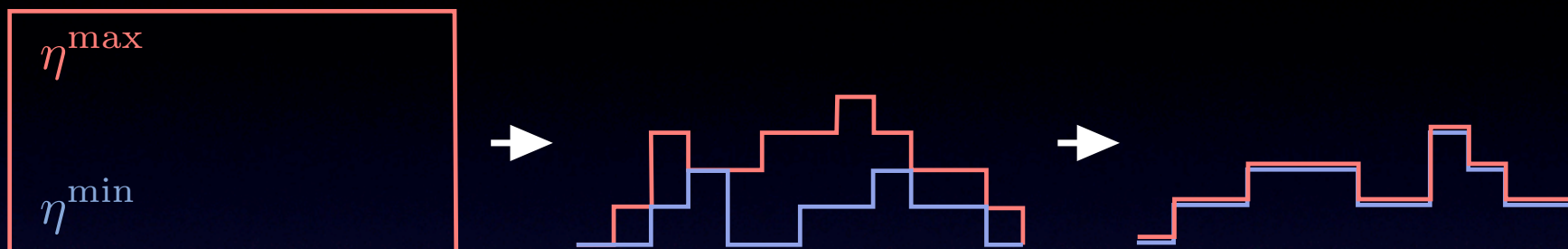
- **Monotone coupling** [Propp/Wilson]



$$E[\Delta \text{ area}] \leq 0 \Rightarrow \text{mixing time} = \tilde{O}(n^5)$$

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$$E[\Delta \text{ area}] \leq 0 \Rightarrow \text{mixing time} = \tilde{O}(n^5)$$

- **Comparison** [Diaconis/Saloff-Coste]

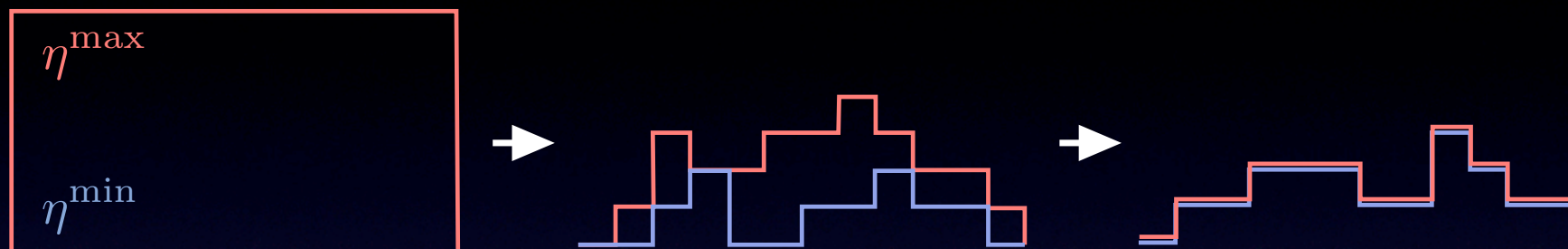
Compare with “non-local” dynamics:

$\eta(i) \rightarrow \text{any value in } [0, n] \text{ w. prob. } \pi(\cdot \mid \eta(i \pm 1))$

$\Rightarrow \text{Mixing time} = \tilde{O}(n^8)$

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Both are loose and give little geometrical insight

Results

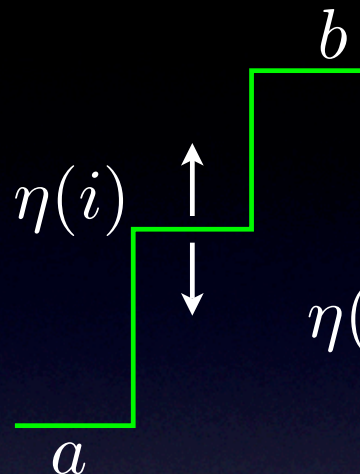
Main Theorem: For all $\beta > 0$, the single-site dynamics has mixing time $\tilde{O}(n^{3.5})$

Also: Almost matching lower bound $\Omega(n^3)$

Bonus: Analysis gives insight into actual evolution of contour

“It’s not a good idea to do a proof in a talk”
[D. Gamarnik, Phys. of Algorithms, 2009]

Non-local “Column” Dynamics



$$\eta(i) \rightarrow k \text{ w. prob. } \propto \begin{cases} \frac{1}{a-b} & a \leq k \leq b \\ e^{-\beta(a-k)} & 0 \leq k < a \\ e^{-\beta(k-b)} & b < k \leq n \end{cases}$$

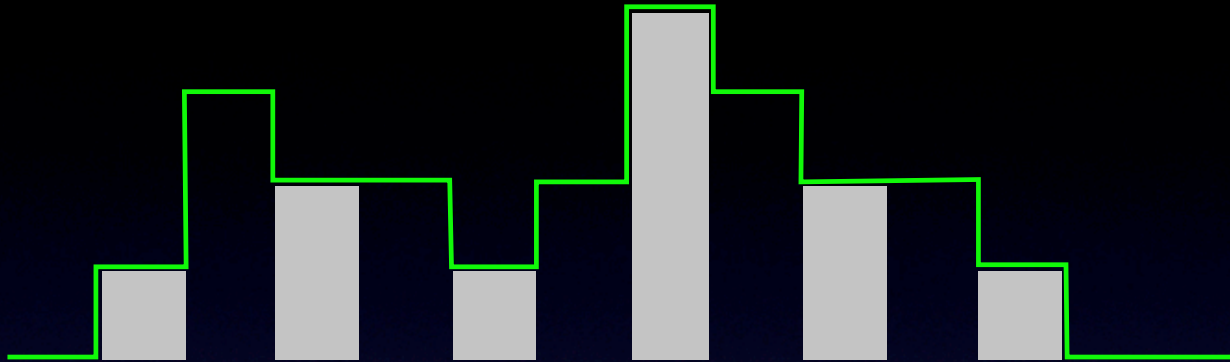
In absence of “walls” at height $0, n$, by symmetry:

$$\mathbb{E}[\eta(i)] = \frac{1}{2}[a + b] = \frac{1}{2}[\eta(i-1) + \eta(i+1)]$$

and mixing time $O(n^3 \log n)$ follows from 2nd eigenvalue of discrete Laplacian

Can show : repulsion from walls only helps!

Simulating Column Dynamics by Single-Site Dynamics



Mixing time of “odd-even” column dynamics is $\tilde{O}(n^2)$

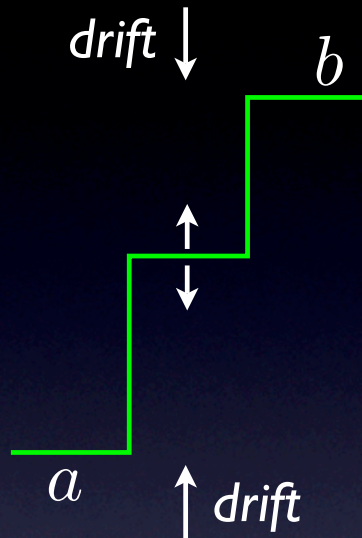
To simulate one move:

- perform $O(t^* n \log n)$ single-site updates, where $t^* =$ mixing time within column
- cancel updates in even (odd) columns

[Peres/Winkler]: Censoring can only slow down dynamics

Hence single-site mixing time is $\tilde{O}(n^3 t^*)$

Mixing time within column, t^*



Mixing time $t^* = \tilde{O}((b - a)^2) \leq \tilde{O}(\text{grad}^2)$
where

$$\text{grad} = \max_i |\eta(i) - \eta(i - 1)|$$

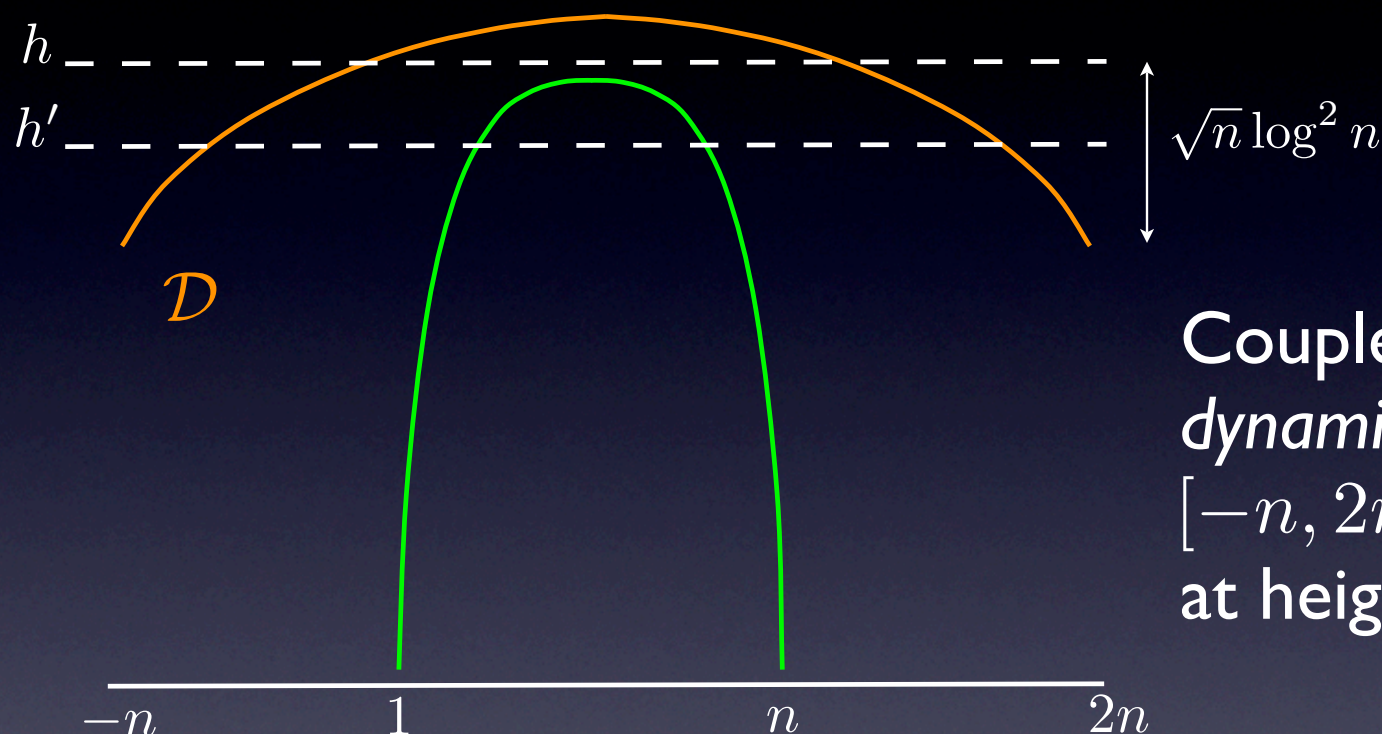
So we need to keep grad small

Problems :

- $\text{grad} = \Theta(n)$ at boundaries !
- Proving non-equilibrium properties for MCs is hard !

Bounding dynamics

Goal: bring contour down from height h to $h' = h - \sqrt{n}$



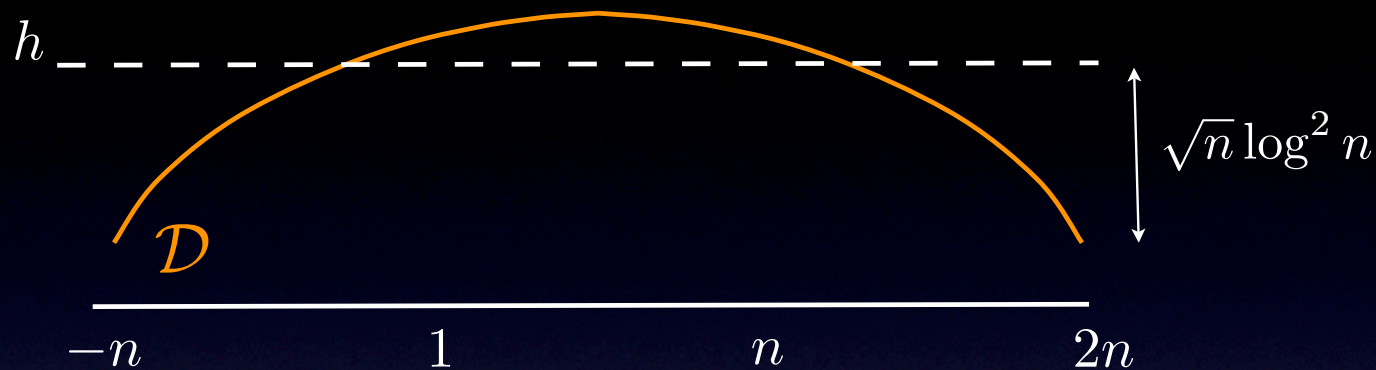
Couple with *bounding dynamics* \mathcal{D} on $[-n, 2n]$ with b.c.'s at height $h - \sqrt{n} \log^2 n$

In equilibrium \mathcal{D} is below h'

Claim: \mathcal{D} has $\text{grad} \leq \text{polylog}(n)$ w.h.p.

Hence: Time to reach small height ($\sim \sqrt{n}$) is $\tilde{O}(n^{3.5})$ ✓

Gradient of bounding dynamics



Event $B = \{\text{grad} > \log^{4.5} n\}$

Easy to see: in equilibrium $\pi(B) \leq e^{-c \log^{4.5} n}$

Start in equilibrium conditional on being above h on $[1, n]$

Note that $\pi(H) \geq e^{-c \log^4 n}$

Thus at all times t , $\Pr_t(B) \leq \pi(B)/\pi(H) \leq e^{-c \log^{4.5} n}$ ✓

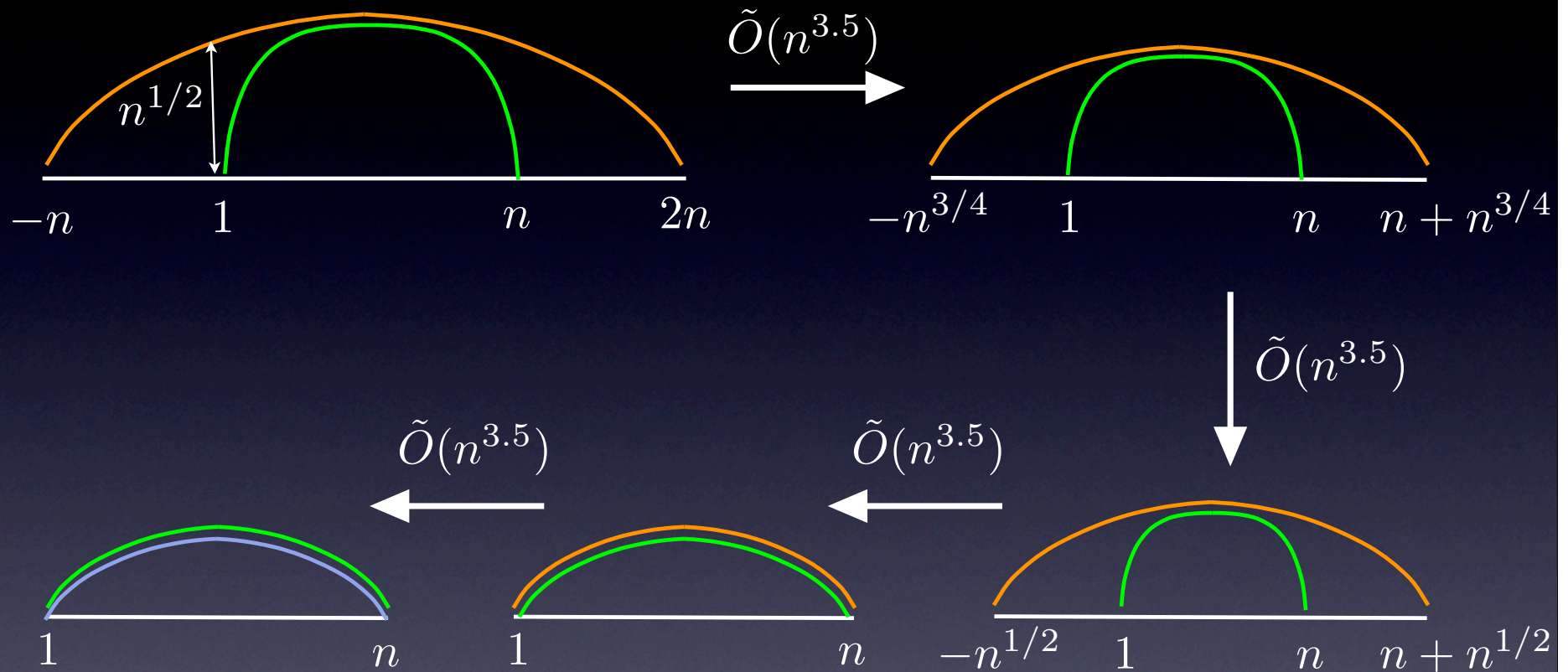
Gradient of bounding dynamics (cont.)

$$B = \{\text{grad} > \log^{4.5} n\} \quad (\text{bad event})$$

$$H = \{\text{above } h \text{ on } [1, n]\} \quad (\text{conditioning event})$$

$$\begin{aligned} \Pr_t(B) &= \sum_{\eta \in B} \sum_{\eta_0} \Pr(\eta_0) \Pr(\eta_t = \eta \mid \eta_0) \\ &= \sum_{\eta \in B} \sum_{\eta_0} \frac{\pi(\eta_0)}{\pi(H)} \Pr(\eta_t = \eta \mid \eta_0) \\ &= \frac{1}{\pi(H)} \sum_{\eta \in B} \pi(\eta) \\ &= \frac{\pi(B)}{\pi(H)} \end{aligned}$$

Final smoothing



So overall mixing time is $\tilde{O}(n^{3.5})$ ✓

Note: Smoothing phase can be improved to $\tilde{O}(n^3)$

Extensions?

- Match the lower bound of $\Omega(n^3)$
- Adapt to lozenge tilings [Luby/Randall/S., Wilson]
- Extend to low-temperature Ising model (see very recent developments by [Martinelli/Toninelli])
- Censoring + Geometry \rightarrow analysis of other (monotone) models